

Constraint Handling Rules

A Tutorial for (PROLOG) Programmers

Tom Schrijvers

Katholieke Universiteit Leuven, Belgium
TOM.SCHRIJVERS@CS.KULEUVEN.BE

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- 2 The Art of CHR
- 3 Simpagation
- 4 For the Sake of Arguments
- 5 En Garde
- 6 A Perfect Match
- 7 Bigger Programs
- 8 Propagation
- 9 Order in the Rules
- 10 Reactivation
- 11 CHR vs. Prolog
- 12 Tutorial Summary
- 13 Facts about CHR



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Expected background:

- ▶ know & like programming

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- ▶ know Prolog, ...

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- ▶ know Prolog, ...
- ▶ but have an itch that Prolog doesn't scratch

Expected background:

- ▶ know & like programming
- ▶ know Prolog, ...
- ▶ but have an itch that Prolog doesn't scratch
- ▶ ...or you will have at the end of this tutorial

About This Tutorial

Expected mindset:

- ▶ forget about Prolog for now

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You've been warned!

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You've been warned! Also:

- ▶ do not expect anything from the word “constraint”

Expected mindset:

- ▶ forget about Prolog for now
- ▶ do not assume things to work as in Prolog
- ▶ really, they won't be the same!!!

You've been warned! Also:

- ▶ do not expect anything from the word “constraint”
- ▶ CHR constraints aren't

What you'll learn:

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- 1 how CHR works (operationally)

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- 1 how CHR works (operationally)
- 2 how to program in CHR

What you'll learn:

- 1 how CHR works (operationally)
- 2 how to program in CHR
- 3 just maybe, what CHR is good for



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As novice, you have to start at the bottom of the ladder:

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- ▶ Learning how to obtain different colors

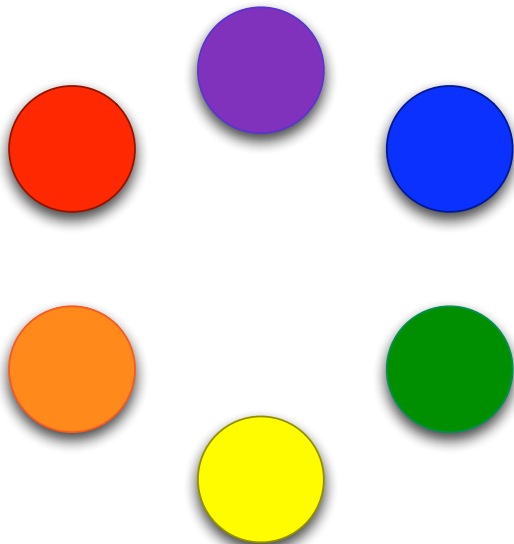
As novice, you have to start at the bottom of the ladder:

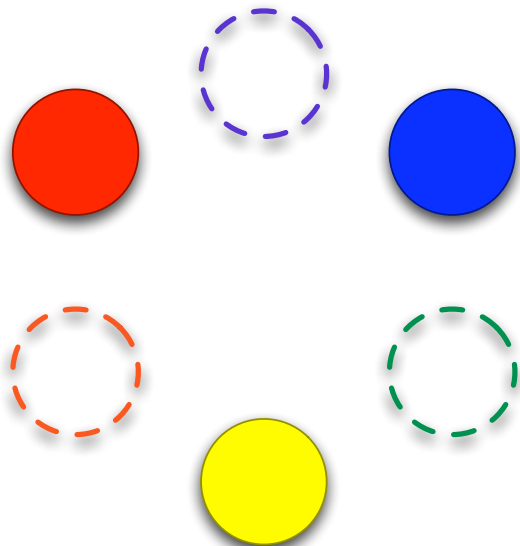
- ▶ Learning how to obtain different colors
- ▶ by mixing paint

As novice, you have to start at the bottom of the ladder:

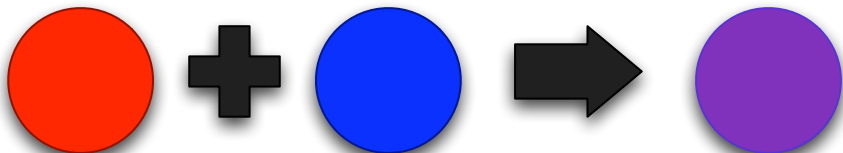
- ▶ Learning how to obtain different colors
- ▶ by mixing paint
- ▶ Made easy by the Paint Mixing Rules (PMR) language

The Paint

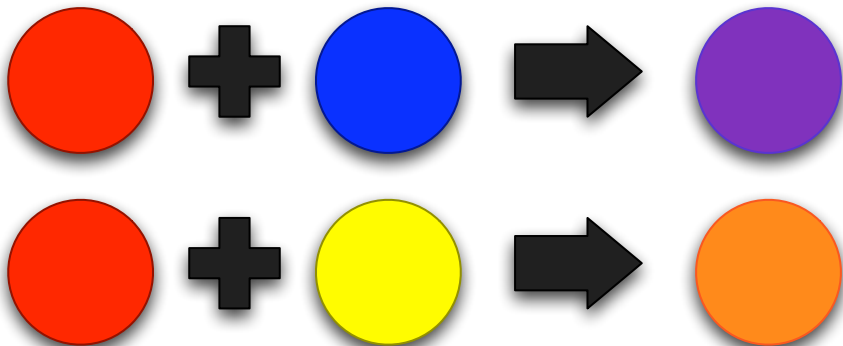




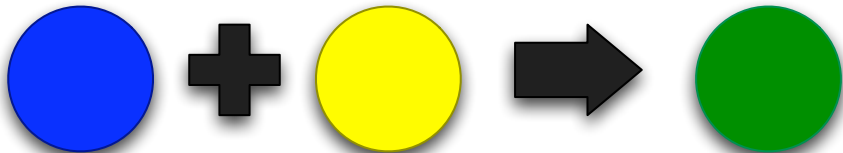
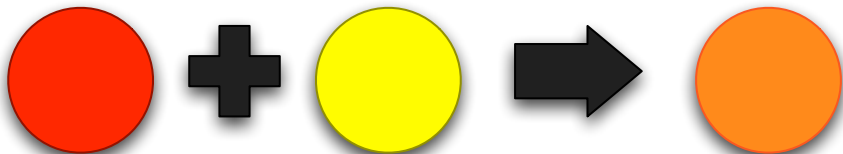
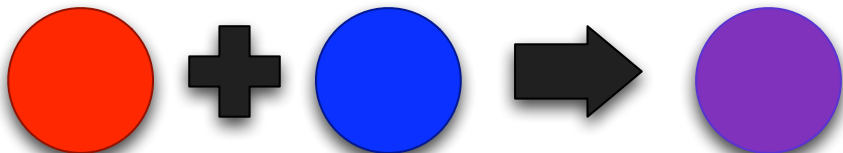
The Rules



The Rules



The Rules



Under the hood, PMR is of course CHR:

PMR	CHR
paints	constraints
rules	simplification rules (textual)
mixing bucket	constraint store

Declaring our paint colors, in textual form:

```
:- chr_constraint red.
```

%



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```
:- chr_constraint red.
```

%



```
:- chr_constraint blue.
```

%



Declaring our paint colors, in textual form:

```
:- chr_constraint red.
```

%



```
:- chr_constraint blue.
```

%



```
:- chr_constraint yellow.
```

%



```
:- chr_constraint orange.
```

%



```
...
```

Are paint colors really “constraints”?

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No: not in the Constraint Programming sense

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Yes: in the CHR sense

- ▶ 1st-class entities in CHR
- ▶ rewritten by rules

Are paint colors really “constraints”?

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Yes: in the CHR sense

- ▶ 1st-class entities in CHR
- ▶ rewritten by rules

(Are Prolog predicates really “predicates”?)

Simplification Rule: $head \Leftrightarrow body$.

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“ $head$ may be replaced with $body$ ”

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red, blue \Leftrightarrow purple.

Simplification Rule: $head \Leftrightarrow body$.

“ $head$ may be replaced with $body$ ”



red, blue \Leftrightarrow purple.



red, yellow \Leftrightarrow orange.

Simplification Rule: $head \Leftrightarrow body$.

“ $head$ may be replaced with $body$ ”

%  +  \Rightarrow 

red, blue \Leftrightarrow purple.

%  +  \Rightarrow 

red, yellow \Leftrightarrow orange.

%  +  \Rightarrow 

blue, yellow \Leftrightarrow green.

Time For...



... a cocktail break!

You can have as many heads as you like:

```
vodka, martini <=> vodka_martini.
```

You can have as many heads as you like:

```
vodka, martini <=> vodka_martini.
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```
tequila, cointreau, lime <=> margarita.
```

You can have as many heads as you like:

```
vodka, martini <=> vodka_martini.
```

```
tequila, cointreau, lime <=> margarita.
```

```
gin, cherry_brandy, cointreau, grenadine,  
pineapple, lemon, angostura_bitters  
  <=> singapore_sling.
```

but at least one!

Warning: too many heads may cause a headache!

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```
vodka_martini, margarita, singapore_sling  
=> hangover, blackout.
```

Warning: too many heads may cause a headache!

```
vodka_martini, margarita, singapore_sling  
=> hangover, blackout.
```

You can have many constraints in the body too.

Paint-Mixing Engine:

- ▶ Prolog (SWI-Prolog, ...)
- ▶ Leuven CHR system

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Everything in a file:

```
paint.pl
:- use_module(library(chr)).

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...

red, blue <=> purple.
...
```

Paint-Mixing Engine:

- ▶ Prolog (SWI-Prolog, ...)
- ▶ Leuven CHR system

Everything in a file:

```
paint.pl
:- use_module(library(chr)).

:- chr_constraint red.
...

red, blue <=> purple.
...
```



```
?- [paint].
```

Running the Rules

```
red, blue    <=> purple.  
red, yellow  <=> orange.  
blue, yellow <=> green.
```

```
?- red, blue.
```

- ▶ queries
- ▶ answers

Running the Rules

```
red, blue    <=> purple.  
red, yellow  <=> orange.  
blue, yellow <=> green.
```

- ▶ queries
- ▶ answers

```
?- red, blue.
```

```
purple
```

```
?- red, yellow.
```

Running the Rules

```
red, blue    <=> purple.  
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blue, yellow <=> green.
```

- ▶ queries
- ▶ answers

```
?- red, blue.
```

```
purple
```

```
?- red, yellow.
```

```
orange
```

```
?- blue, yellow.
```

Running the Rules

```
red, blue    <=> purple.  
red, yellow  <=> orange.  
blue, yellow <=> green.
```

- ▶ queries
- ▶ answers

```
?- red, blue.
```

```
purple
```

```
?- red, yellow.
```

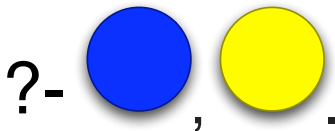
```
orange
```

```
?- blue, yellow.
```

```
green
```

In Slow Motion

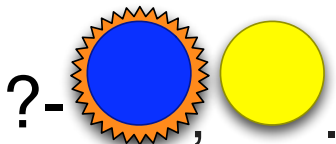
```
red, blue    <=> purple.  
red, yellow  <=> orange.  
blue, yellow <=> green.
```



▶ query:
left-to-right

In Slow Motion

```
red, blue    <=> purple.  
red, yellow  <=> orange.  
blue, yellow <=> green.
```

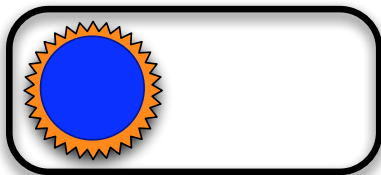


- ▶ query:
left-to-right
- ▶ active
constraint:
blue

In Slow Motion

```
red, blue    <=> purple.  
red, yellow  <=> orange.  
blue, yellow <=> green.
```

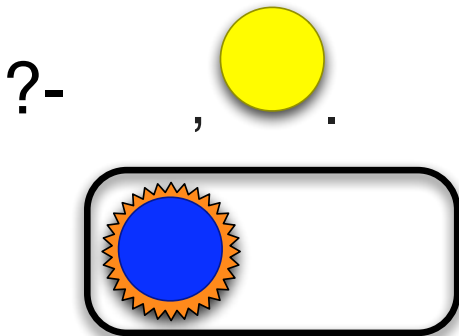
?-



- ▶ query:
left-to-right
- ▶ active
constraint:
blue
- ▶ put in
constraint
store

In Slow Motion

```
red, blue    <=> purple.  
red, yellow  <=> orange.  
blue, yellow <=> green.
```

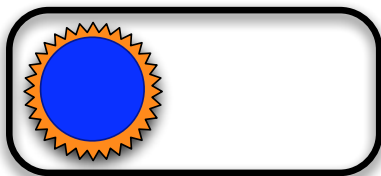


- ▶ active constraint: blue
- ▶ fire first rule?
- ▶ no partner red in constraint store

In Slow Motion

```
red, blue    <=> purple.  
red, yellow  <=> orange.  
blue, yellow <=> green.
```

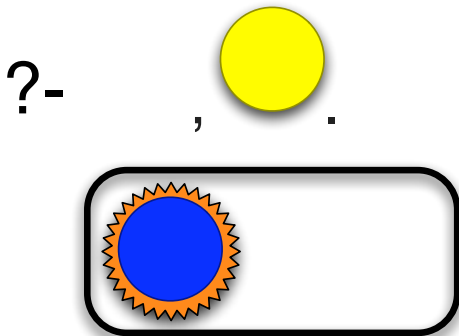
?-



- ▶ active constraint: blue
- ▶ fire second rule?
- ▶ no blue in head

In Slow Motion

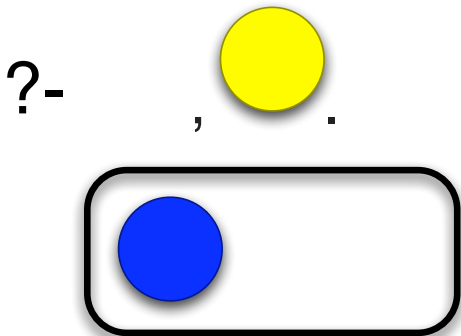
```
red, blue    <=> purple.  
red, yellow  <=> orange.  
blue, yellow <=> green.
```



- ▶ active constraint: blue
- ▶ fire third rule?
- ▶ no partner yellow in constraint store

In Slow Motion

```
red, blue    <=> purple.  
red, yellow  <=> orange.  
blue, yellow <=> green.
```

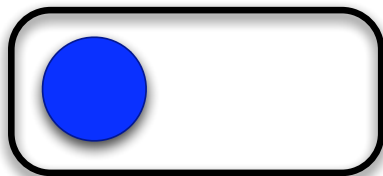
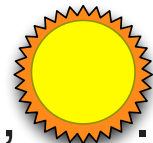


- ▶ deactivate constraint blue

In Slow Motion

```
red, blue    <=> purple.  
red, yellow  <=> orange.  
blue, yellow <=> green.
```

?-



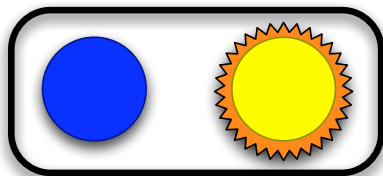
▶ active
constraint:
yellow

In Slow Motion

```
red, blue    <=> purple.  
red, yellow  <=> orange.  
blue, yellow <=> green.
```

?-

,



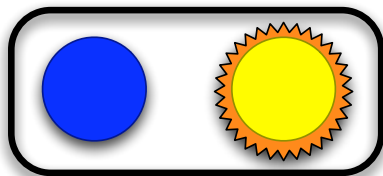
- ▶ active constraint: yellow
- ▶ put in constraint store

In Slow Motion

```
red, blue    <=> purple.  
red, yellow  <=> orange.  
blue, yellow <=> green.
```

?-

, .



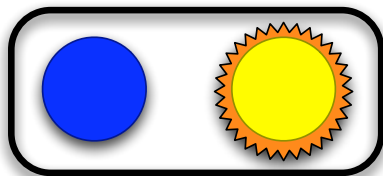
- ▶ active constraint: yellow
- ▶ fire first rule?
- ▶ no yellow in head

In Slow Motion

```
red, blue    <=> purple.  
red, yellow  <=> orange.  
blue, yellow <=> green.
```

?-

,



- ▶ active constraint: yellow
- ▶ fire second rule?
- ▶ no partner red in constraint store

In Slow Motion

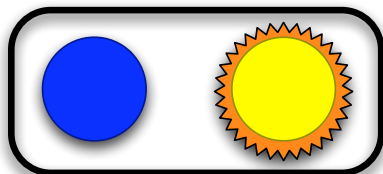
red, blue \Leftrightarrow purple.

red, yellow \Leftrightarrow orange.

blue, yellow \Leftrightarrow green.

?-

,



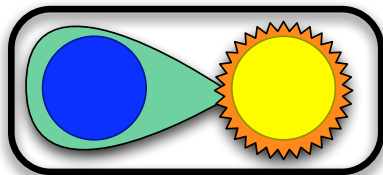
- ▶ active constraint: yellow
- ▶ fire third rule?

In Slow Motion

```
red, blue    <=> purple.  
red, yellow  <=> orange.  
blue, yellow <=> green.
```

?-

,



- ▶ active constraint: yellow
- ▶ fire third rule!
- ▶ with partner blue in constraint store

In Slow Motion

```
red, blue    <=> purple.  
red, yellow  <=> orange.  
blue, yellow <=> green.
```

?-

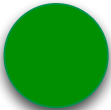
, .



- ▶ fire third rule!
- ▶ delete the
matched
head
constraints

In Slow Motion

```
red, blue    <=> purple.  
red, yellow  <=> orange.  
blue, yellow <=> green.
```

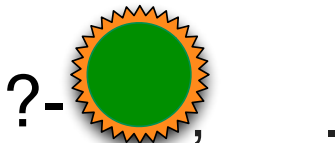
?- , .



- ▶ fire third rule!
- ▶ add the body to (front of) query

In Slow Motion

```
red, blue    <=> purple.  
red, yellow <=> orange.  
blue, yellow <=> green.
```



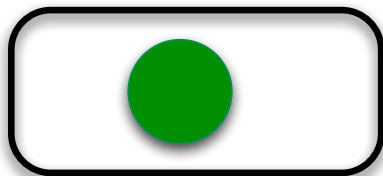
- ▶ active constraint: green

In Slow Motion

```
red, blue    <=> purple.  
red, yellow  <=> orange.  
blue, yellow <=> green.
```

?-

, .



to make a long
story short:

- ▶ ...
- ▶ final
constraint
store

What Happens?

```
red, blue    <=> purple.  
red, yellow  <=> orange.  
blue, yellow <=> green.
```

```
?- red, yellow.  
orange
```


What Happens?

```
red, blue    <=> purple.  
red, yellow <=> orange.  
blue, yellow <=> green.
```

```
?- red, yellow.
```

```
orange
```

```
?- yellow, red.
```

```
orange
```

What Happens?

```
red, blue    <=> purple.  
red, yellow <=> orange.  
blue, yellow <=> green.
```

```
?- red.
```

What Happens?

```
red, blue    <=> purple.  
red, yellow <=> orange.  
blue, yellow <=> green.
```

```
?- red.
```

```
red
```

What Happens?

```
red, blue    <=> purple.  
red, yellow <=> orange.  
blue, yellow <=> green.
```

```
?- red, yellow, blue.
```

What Happens?

```
red, blue    <=> purple.  
red, yellow  <=> orange.  
blue, yellow <=> green.
```

```
?- red, yellow, blue.
```

```
orange
```

```
blue
```

What Happens?

```
red, blue    <=> purple.  
red, yellow  <=> orange.  
blue, yellow <=> green.
```

```
?- red, blue, yellow.
```

What Happens?

```
red, blue    <=> purple.  
red, yellow  <=> orange.  
blue, yellow <=> green.
```

```
?- red, blue, yellow.
```

```
purple
```

```
yellow
```

Simplification Rule:

- ▶ $head \Leftrightarrow body$.
- ▶ replace head with body

Simplification Rule:

- ▶ *head* \Leftrightarrow *body* .
- ▶ replace head with body

Queries:

- ▶ processed left-to-right
- ▶ incrementally
- ▶ active constraint looks for partners



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Brown stays brown.

```
brown, orange <=> brown.
```

```
brown, green <=> brown.
```

```
brown, purple <=> brown.
```

```
...
```

Simpagation rule: $head_k \setminus head_r \Leftrightarrow body.$

```
brown \ orange => true.
```

```
brown \ green  => true.
```

```
brown \ purple => true.
```

```
...
```

Simpagation rule: $head_k \setminus head_r \Leftrightarrow body$.

```
brown \ orange => true.
```

```
brown \ green  => true.
```

```
brown \ purple => true.
```

```
...
```

- ▶ $head_k$: kept head (1 or more)

Simpagation rule: $head_k \setminus head_r \Leftrightarrow body$.

```
brown \ orange => true.
```

```
brown \ green  => true.
```

```
brown \ purple => true.
```

```
...
```

- ▶ $head_k$: kept head (1 or more)
- ▶ $head_r$: removed head (1 or more)

Simpagation rule: $head_k \setminus head_r \Leftrightarrow body$.

```
brown \ orange => true.
```

```
brown \ green  => true.
```

```
brown \ purple => true.
```

```
...
```

- ▶ $head_k$: kept head (1 or more)
- ▶ $head_r$: removed head (1 or more)
- ▶ true : Prolog's no-op

```
philosophers_stone \ lead  
=> gold.
```

```
?- philosophers_stone, lead.
```



```
philosophers_stone \ lead  
=> gold.
```

```
?- philosophers_stone, lead.  
philosophers_stone  
gold
```

```
philosophers_stone \ lead  
=> gold.
```

```
?- philosophers_stone, lead.
```

```
philosophers_stone
```

```
gold
```

```
?- philosophers_stone, lead, lead.
```

```
philosophers_stone \ lead  
=> gold.
```

```
?- philosophers_stone, lead.
```

```
philosophers_stone
```

```
gold
```

```
?- philosophers_stone, lead, lead.
```

```
philosophers_stone
```

```
gold
```

```
gold
```

All Combinations

```
philosophers_stone \ lead  
=> gold.
```

```
?- lead, lead.
```

All Combinations

```
philosophers_stone \ lead  
=> gold.
```

```
?- lead, lead.
```

```
lead
```

```
lead
```

All Combinations

```
philosophers_stone \ lead  
=> gold.
```

```
?- lead, lead.
```

```
lead
```

```
lead
```

```
?- lead, lead, philosophers_stone.
```

All Combinations

```
philosophers_stone \ lead  
=> gold.
```

```
?- lead, lead.
```

```
lead
```

```
lead
```

```
?- lead, lead, philosophers_stone.
```

```
philosophers_stone
```

```
gold
```

```
gold
```

```
philosophers_stone \ lead  
=> gold.
```

```
?- lead, lead.
```

```
lead
```

```
lead
```

```
?- lead, lead, philosophers_stone.
```

```
philosophers_stone
```

```
gold
```

```
gold
```

Rule fires with all combinations!

Simplification Rule:

- ▶ $head \Leftrightarrow body$.
- ▶ replace head with body

Simplification Rule:

- ▶ $head \Leftrightarrow body$.
- ▶ replace head with body

Simpagation Rule:

- ▶ $head_k \setminus head_r \Leftrightarrow body$.
- ▶ replace $head_r$ with $body$
- ▶ in the presence of $head_k$
- ▶ fires all possible combinations



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```
:- chr_constraint piggy/1.
```

```
piggy(I), piggy(J) <=> K is I + J, piggy(K).
```

- ▶ constraints can have arguments

```
:- chr_constraint piggy/1.
```

```
piggy(I), piggy(J) <=> K is I + J, piggy(K).
```

- ▶ constraints can have arguments
- ▶ arity must be declared

```
:- chr_constraint piggy/1.
```

```
piggy(I), piggy(J) <=> K is I + J, piggy(K).
```

- ▶ constraints can have arguments
- ▶ arity must be declared
- ▶ Prolog terms as arguments

Piggy Bank Merger

```
:- chr_constraint piggy/1.
```

```
piggy(I), piggy(J) <=> K is I + J, piggy(K).
```

?- piggy(5), piggy(1), piggy(4), piggy(2).



Piggy Bank Merger

```
:- chr_constraint piggy/1.
```




```
piggy(I), piggy(J) <=> K is I + J, piggy(K).
```



Piggy Bank Merger

```
:- chr_constraint piggy/1.
```

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?- ,  ,  ,  .

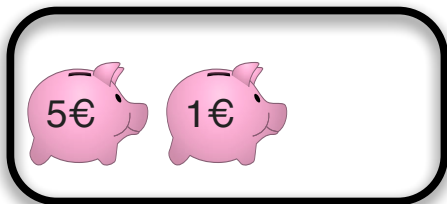


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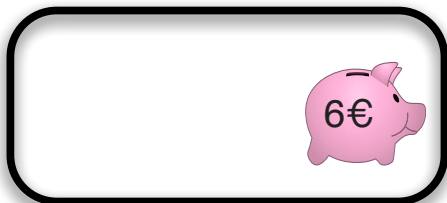


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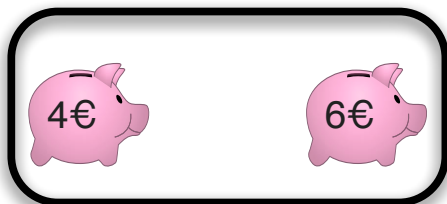


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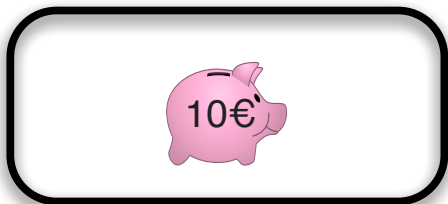


Piggy Bank Merger

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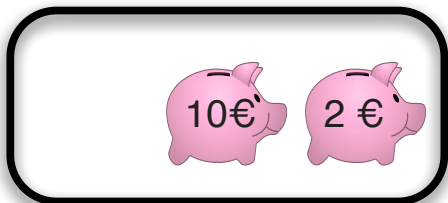
Piggy Bank Merger

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, , , .



Piggy Bank Merger

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?-

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CHR is an embedded language

- ▶ embedded in host language \mathcal{X}

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- ▶ here $\mathcal{X} = \text{Prolog}$

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2-way communication:

CHR is an embedded language

- ▶ embedded in host language \mathcal{X}
- ▶ here $\mathcal{X} = \text{Prolog}$

2-way communication:

- ▶ Prolog calls CHR constraints
e.g. all along from toplevel

CHR is an embedded language

- ▶ embedded in host language \mathcal{X}
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2-way communication:

- ▶ Prolog calls CHR constraints
e.g. all along from toplevel
- ▶ CHR calls Prolog
e.g. `K is I + J`

Not Just Numbers

```
:- chr_constraint value/1.
```

```
value(I), value(J) <=> append(I,J,K), value(K).
```

Not Just Numbers

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:- chr_constraint value/1.
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```
value(I), value(J) <=> append(I,J,K), value(K).
```

```
?- value([a]), value([b]).
```

Not Just Numbers

```
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value(I), value(J) <=> append(I,J,K), value(K).
```

```
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value([b,a]).
```

Not Just Numbers

```
:- chr_constraint value/1.
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value(I), value(J) <=> append(I,J,K), value(K).
```

```
?- value([a]), value([b]).  
value([b,a]).
```

```
?- value([f(a),g(b)]), value([h(i,j)]).
```


Not Just Numbers

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value([b,a]).
```

```
?- value([f(a),g(b)]), value([h(i,j)]).  
value([h(i,j),f(a),g(b)]).
```

Not Just Numbers

```
:- chr_constraint value/1.
```

```
value(I), value(J) <=> append(I,J,K), value(K).
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```
?- value([a]), value([b]).  
value([b,a]).
```

```
?- value([f(a),g(b)]), value([h(i,j)]).  
value([h(i,j),f(a),g(b)]).
```

Any Prolog terms!

CHR constraints

- ▶ zero or more arguments
- ▶ Prolog terms

CHR constraints

- ▶ zero or more arguments
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Bodies

- ▶ CHR constraints and
- ▶ Prolog predicate calls



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Problem: write a program to enumerate values:

```
?- generate(5).
```

Problem: write a program to enumerate values:

```
?- generate(5).  
value(5)  
value(4)  
value(3)  
value(2)  
value(1)
```

Solution:

```
:- chr_constraint generate/1.
```


Solution:

```
:- chr_constraint generate/1.  
:- chr_constraint value/1.
```

Solution:

```
:- chr_constraint generate/1.
```

```
:- chr_constraint value/1.
```

```
generate(N) <=>
```

```
value(N), M is N - 1, generate(M).
```

Running the Program

```
:- chr_constraint generate/1.
```

```
:- chr_constraint value/1.
```

```
generate(N) <=> value(N), M is N - 1, generate(M).
```

```
?- generate(5).
```

Running the Program

```
:- chr_constraint generate/1.  
:- chr_constraint value/1.
```

```
generate(N) <=> value(N), M is N - 1, generate(M).
```

```
?- generate(5).  
    % waiting
```

Running the Program

```
:- chr_constraint generate/1.  
:- chr_constraint value/1.
```

```
generate(N) <=> value(N), M is N - 1, generate(M).
```

```
?- generate(5).  
    % waiting  
    % still waiting
```

Running the Program

```
:- chr_constraint generate/1.  
:- chr_constraint value/1.
```

```
generate(N) <=> value(N), M is N - 1, generate(M).
```

```
?- generate(5).  
    % waiting  
    % still waiting  
ERROR: Out of local stack
```

Running the Program

```
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generate(N) <=> value(N), M is N - 1, generate(M).
```

```
?- generate(5).  
    % waiting  
    % still waiting  
ERROR: Out of local stack
```

Recursion without a base case!

Guarded Rules

```
:- chr_constraint generate/1.
```

```
:- chr_constraint value/1.
```

```
generate(N) <=> N == 0 | true.
```

```
generate(N) <=> N > 0 | value(N),  
                M is N - 1, generate(M).
```



```
:- chr_constraint generate/1.  
:- chr_constraint value/1.  
  
generate(N) <=> N == 0 | true.  
generate(N) <=> N > 0 | value(N),  
                    M is N - 1, generate(M).
```

Guarded rule:

- ▶ apply rule **if** guard succeeds

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```

Guarded rule:

- ▶ apply rule **if** guard succeeds
- ▶ guard is optional
- ▶ guard may contain any Prolog, but may not bind any variables from the head (pure check)

Even Shorter

```
generate(N) <=> N > 0 | value(N),  
                M is N - 1, generate(M).
```

Even Shorter

```
generate(N) <=> N > 0 | value(N),  
                M is N - 1, generate(M).
```

```
?- generate(5).
```

Even Shorter

```
generate(N) <=> N > 0 | value(N),  
                M is N - 1, generate(M).
```

```
?- generate(5).  
value(5)  
value(4)  
value(3)  
value(2)  
value(1)  
generate(0)
```

Simplification Rule:

- ▶ $head \Leftrightarrow guard \mid body .$
- ▶ replace head with body
- ▶ if $guard$ succeeds

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Simpagation Rule:

- ▶ $head_k \setminus head_r \Leftrightarrow guard \mid body .$
- ▶ replace $head_r$ with $body$
- ▶ in the presence of $head_k$
- ▶ if $guard$ succeeds



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generate(0) <=> true.  
generate(N) <=> N > 0 | value(N),  
                    M is N - 1, generate(M).
```

Matching:

- ▶ “inline” notation for guard
- ▶ same meaning as explicit guard

Meaning of Matching

Maching	Guard
$c(X) \Leftrightarrow \text{true}.$	$c(X) \Leftrightarrow \text{true} \mid \text{true}.$

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$c(X) \Leftrightarrow \text{true}.$	$c(X) \Leftrightarrow \text{true} \mid \text{true}.$
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Matching	Guard
$c(X) \Leftarrow \text{true}.$	$c(X) \Leftarrow \text{true} \mid \text{true}.$
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$c(X, X) \Leftarrow \text{true}.$	$c(X, Y) \Leftarrow X == Y \mid \text{true}.$

Meaning of Matching

Maching	Guard
$c(X) \langle \Rightarrow \rangle \text{true.}$	$c(X) \langle \Rightarrow \rangle \text{true} \mid \text{true.}$
$c(a) \langle \Rightarrow \rangle \text{true.}$	$c(X) \langle \Rightarrow \rangle X == a \mid \text{true.}$
$c(f(A)) \langle \Rightarrow \rangle \text{true.}$	$c(X) \langle \Rightarrow \rangle \text{nonvar}(X), X = f(A) \mid \text{true.}$
$c(X, X) \langle \Rightarrow \rangle \text{true.}$	$c(X, Y) \langle \Rightarrow \rangle X == Y \mid \text{true.}$
$c(X), d(X) \langle \Rightarrow \rangle \text{true.}$	$c(X), d(Y) \langle \Rightarrow \rangle X == Y \mid \text{true.}$

Matching: CHR vs. Prolog

```
p(world) :- writeln('Hello, World!').
```

```
:- chr_constraint c/1.
```

```
c(world) <=> writeln('Hello, World!').
```


Matching: CHR vs. Prolog

```
p(world) :- writeln('Hello, World!').
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```
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```

```
:- chr_constraint c/1.
```

```
c(world) <=> writeln('Hello, World!').
```

```
?- p(world).  
Hello, World!
```

```
?- p(Free).
```

Matching: CHR vs. Prolog

```
p(world) :- writeln('Hello, World!').
```

```
:- chr_constraint c/1.
```

```
c(world) <=> writeln('Hello, World!').
```

```
?- p(world).  
Hello, World!
```

```
?- p(Free).  
Hello, World!  
Free = world
```

Matching: CHR vs. Prolog

```
p(world) :- writeln('Hello, World!').
```

```
:- chr_constraint c/1.
```

```
c(world) <=> writeln('Hello, World!').
```

```
?- p(world).  
Hello, World!
```

```
?- p(Free).  
Hello, World!  
Free = world
```

```
?- p(hello).
```

Matching: CHR vs. Prolog

```
p(world) :- writeln('Hello, World!').
```

```
:- chr_constraint c/1.
```

```
c(world) <=> writeln('Hello, World!').
```

```
?- p(world).  
Hello, World!
```

```
?- p(Free).  
Hello, World!  
Free = world
```

```
?- p(hello).  
No
```

Matching: CHR vs. Prolog

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p(world) :- writeln('Hello, World!').
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```
:- chr_constraint c/1.
```

```
c(world) <=> writeln('Hello, World!').
```

```
?- p(world).  
Hello, World!
```

```
?- p(Free).  
Hello, World!  
Free = world
```

```
?- p(hello).  
No
```

```
?- c(world).
```

Matching: CHR vs. Prolog

```
p(world) :- writeln('Hello, World!').
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?- p(hello).  
No
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Hello, World!  
Free = world
```

```
?- p(hello).  
No
```

```
?- c(world).  
Hello, World!
```

```
?- c(Free).
```

Matching: CHR vs. Prolog

```
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```

```
c(world) <=> writeln('Hello, World!').
```

```
?- p(world).  
Hello, World!
```

```
?- p(Free).  
Hello, World!  
Free = world
```

```
?- p(hello).  
No
```

```
?- c(world).  
Hello, World!
```

```
?- c(Free).  
c(Free)
```

Matching: CHR vs. Prolog

```
p(world) :- writeln('Hello, World!').
```

```
:- chr_constraint c/1.
```

```
c(world) <=> writeln('Hello, World!').
```

```
?- p(world).  
Hello, World!
```

```
?- p(Free).  
Hello, World!  
Free = world
```

```
?- p(hello).  
No
```

```
?- c(world).  
Hello, World!
```

```
?- c(Free).  
c(Free)
```

```
?- c(hello).
```

Matching: CHR vs. Prolog

```
p(world) :- writeln('Hello, World!').
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```

```
c(world) <=> writeln('Hello, World!').
```

```
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Hello, World!
```

```
?- p(Free).  
Hello, World!  
Free = world
```

```
?- p(hello).  
No
```

```
?- c(world).  
Hello, World!
```

```
?- c(Free).  
c(Free)
```

```
?- c(hello).  
c(hello)
```

Matching

- ▶ short-hand for equality-based guards
- ▶ one-way unification
 - ▶ only instantiates head
 - ▶ does not instantiate constraints



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Combining Rules

```
generate(0) <=> true.  
generate(N) <=> N > 0 | value(N),  
                    M is N - 1, generate(M).  
  
value(I), value(J) <=> K is I + J, value(K).
```

What's the outcome?

```
?- generate(4).
```

Combining Rules

```
generate(0) <=> true.  
generate(N) <=> N > 0 | value(N),  
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```

What's the outcome?

```
?- generate(4).  
value(10)
```


Combining Rules

```
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generate(N) <=> N > 0 | value(N),  
                    M is N - 1, generate(M).  
  
value(I), value(J) <=> K is I + J, value(K).
```

What's the outcome?

```
?- generate(4).  
value(10)  
  
?- generate(5).
```

Combining Rules

```
generate(0) <=> true.  
generate(N) <=> N > 0 | value(N),  
                    M is N - 1, generate(M).  
  
value(I), value(J) <=> K is I + J, value(K).
```


What's the outcome?

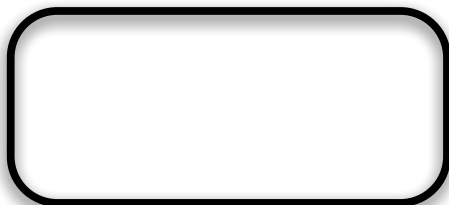
```
?- generate(4).  
value(10)  
  
?- generate(5).  
value(15)
```

In Slow Motion

```
generate(0) <=> true.  
generate(N) <=> N > 0 | value(N),  
                    M is N - 1, generate(M).
```

```
value(I), value(J) <=> K is I + J, value(K).
```

?-  2

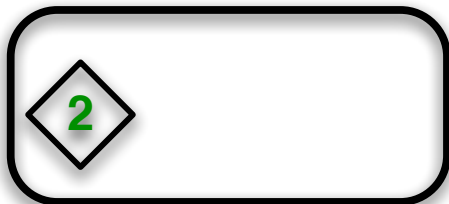


In Slow Motion

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```

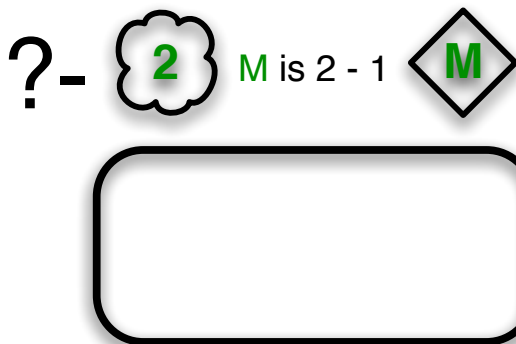
?-



In Slow Motion

```
generate(0) <=> true.  
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```


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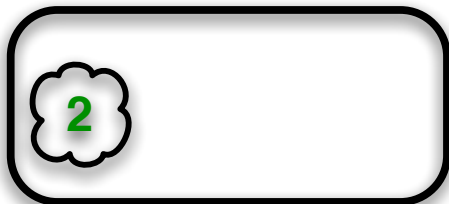


In Slow Motion

```
generate(0) <=> true.  
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```

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value(I), value(J) <=> K is I + J, value(K).
```

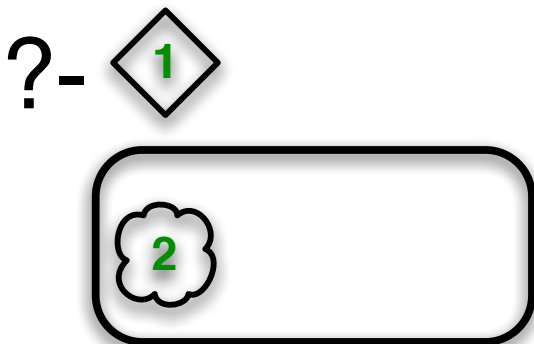
? - M is 2 - 1 



In Slow Motion

```
generate(0) <=> true.  
generate(N) <=> N > 0 | value(N),  
                    M is N - 1, generate(M).
```

```
value(I), value(J) <=> K is I + J, value(K).
```

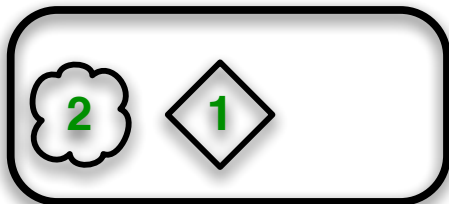


In Slow Motion

```
generate(0) <=> true.  
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value(I), value(J) <=> K is I + J, value(K).
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?-



In Slow Motion

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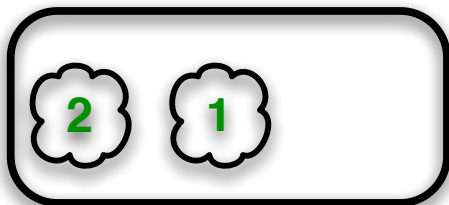
In Slow Motion

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generate(0) <=> true.  
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```

?-

M is 1 - 1



In Slow Motion

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?-

M is 1 - 1

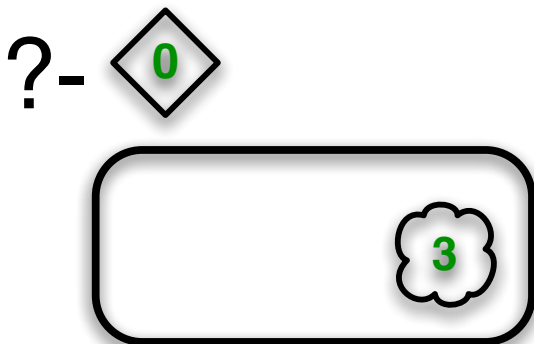
M

3

In Slow Motion

```
generate(0) <=> true.  
generate(N) <=> N > 0 | value(N),  
                    M is N - 1, generate(M).
```

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value(I), value(J) <=> K is I + J, value(K).
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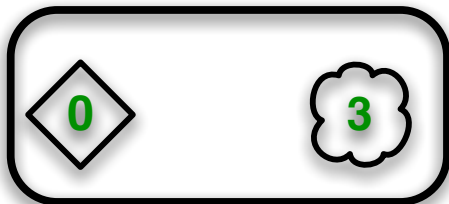


In Slow Motion

```
generate(0) <=> true.  
generate(N) <=> N > 0 | value(N),  
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?-

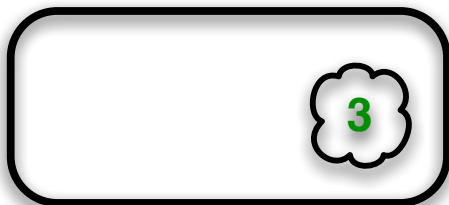


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Combining Rules

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generate(1) <=> true.  
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```

What's the outcome?

```
?- generate(4).
```

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```
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value(2)  
value(3)
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```

```
value(2)
```

```
value(3)
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```
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Just added value(3).
What now?

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► remove value(6) ?

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Just added value(3).
What now?

- ▶ remove value(6) ?
- ▶ remove value(9) ?
- ▶ unspecified in CHR!

Alternate Ending

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generate(1) <=> true.  
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CHR Programs

- ▶ Programs consist of a sequence of rules

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- ▶ Active constraint traverses the rules top-to-bottom to find any that fire
- ▶ Rest of query/body waits
- ▶ Unspecified which partner is found



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Simpagation rule

$head_k \setminus head_r \Leftrightarrow guard \mid body.$

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What if $head_k = \emptyset$?

Simpagation with $H_k = \emptyset$

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Simpagation rule

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What if $head_r = \emptyset$?

Simpagation rule

$head_k \setminus head_r \iff guard \mid body .$

What if $head_r = \emptyset$?

Propagation rule

$head_k \implies guard \mid body .$

- ▶ add $body$
- ▶ in the presence of $head_k$
- ▶ if $guard$ holds

Propagation

$a, b \implies c.$

$?- a, b.$

Propagation

$a, b \implies c.$

?- a, b.

a

b

c

Propagation

$a, b \implies c.$

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a

b

c

?- a, b, b.

Propagation

$a, b \implies c.$

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a

b

c

?- a, b, b.

a

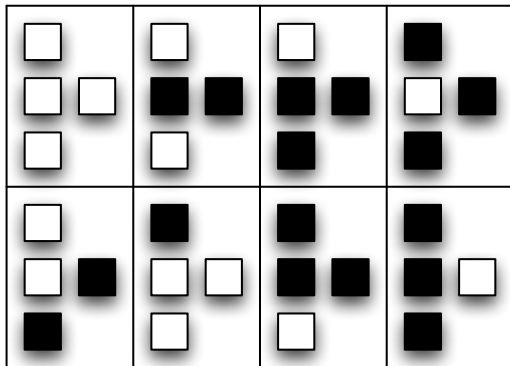
b

b

c

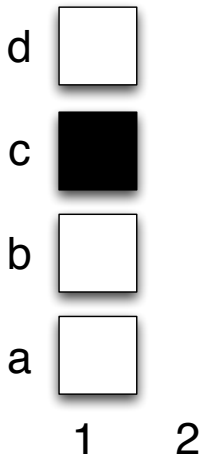
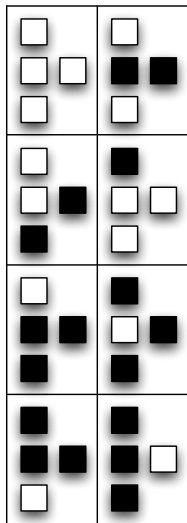
c

Writing Wolfram's cellular automaton Rule 110 in CHR.



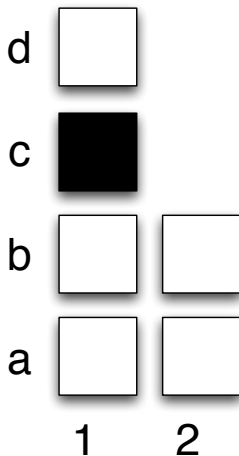
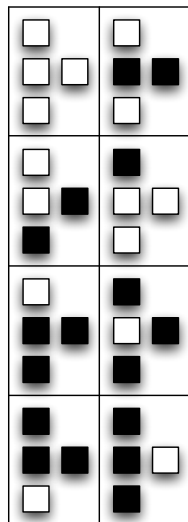
Cellular Automaton

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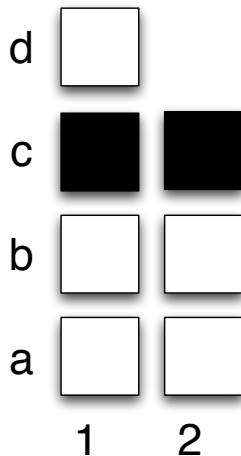
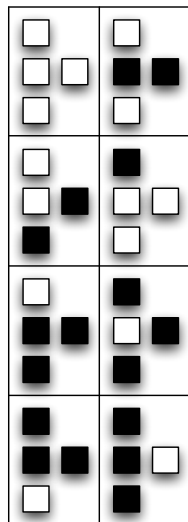
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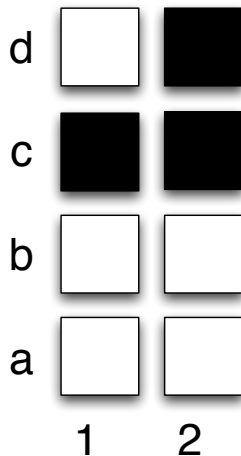
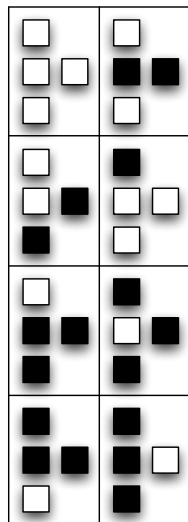
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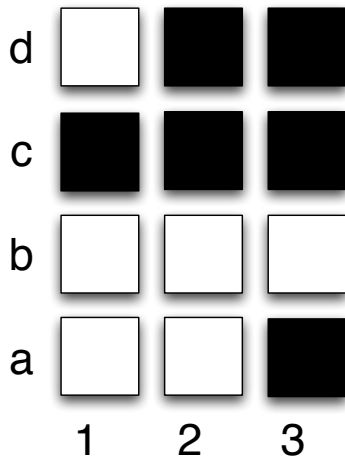
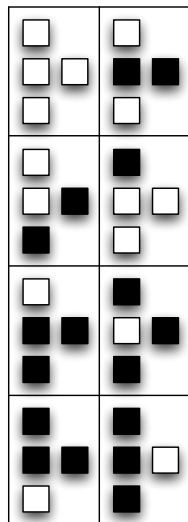
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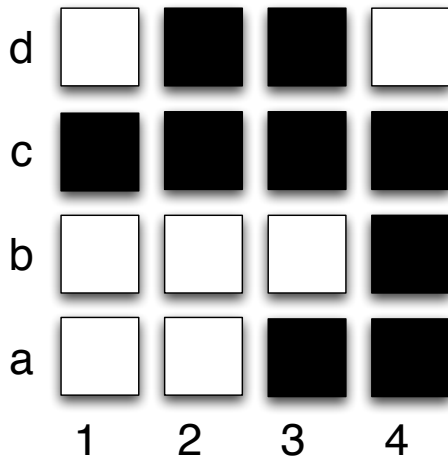
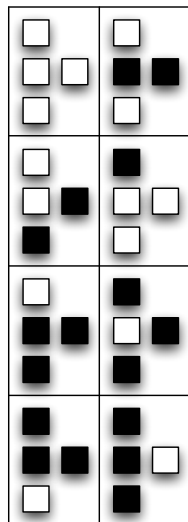
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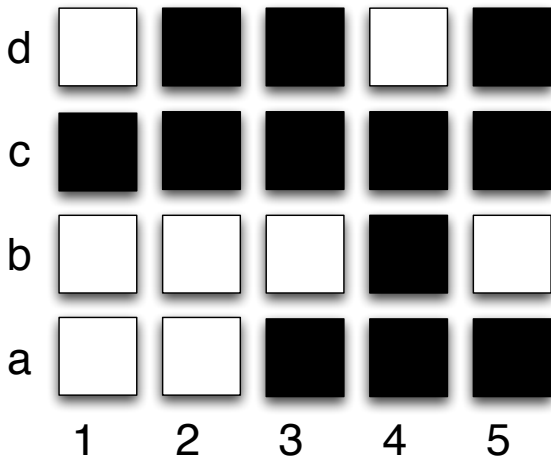
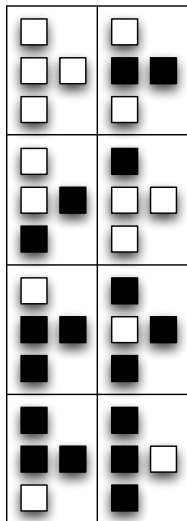
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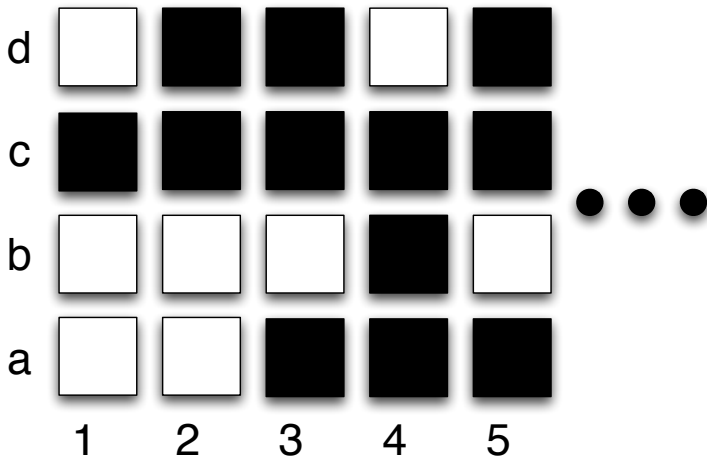
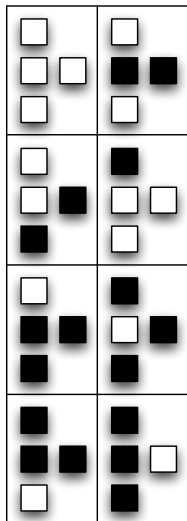
Cellular Automaton

Writing Wolfram's cellular automaton Rule 110 in CHR.



Cellular Automaton

Writing Wolfram's cellular automaton Rule 110 in CHR.



CHR cells

```
% black(Name,Below,Above,Generation).
```

```
:- chr_constraint black/4.
```

```
% white(Name,Below,Above,Generation).
```

```
:- chr_constraint white/4.
```

CHR cells

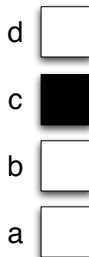
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```
:- chr_constraint black/4.
```

```
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```

```
:- chr_constraint white/4.
```

```
?- white(d,c,a,1),  
    black(c,b,d,1),  
    white(b,a,c,1),  
    white(a,d,b,1).
```



Automaton Rule

```
white(C,_,_.G) ,  
black(B,A,C,G) ,  
white(A,_,_,G)  
==> G < 10 | NG is G + 1, black(B,A,C,NG).
```

Automaton Rule

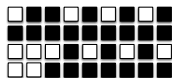
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   white(b,a,c,1),  
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Another Example of Propagation

```
generate(N) ==> value(2).
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generate(N) ==> value(2).
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```
generate(N), value(I) ==>  
    I < N | J is I + 1, value(J).
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```
value(I) \ value(J) <=> J mod I == 0 | true.
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?- generate(10).
```

```
value(2)
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value(3)
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```
value(5)
```

```
value(7)
```

```
generate(10)
```

Simplification Rule: $head \Leftrightarrow guard \mid body .$

- ▶ replace head with body

Simpagation Rule: $head_k \setminus head_r \Leftrightarrow guard \mid body .$

- ▶ replace $head_r$ with $body$
- ▶ in the presence of $head_k$

Propagation Rule: $head \Rightarrow guard \mid body .$

- ▶ add $body$
- ▶ in the presence of $head$
- ▶ fires once for each combination (propagation history)



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Biased Coin?

```
coin <=> heads.  
coin <=> tails.
```

What's the outcome?

```
?- coin.
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Rules are tried in order!

Coin Flipping in Prolog

CHR

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coin <=> heads.  
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```

Prolog

```
p_coin :- heads.  
p_coin :- tails.
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What's the outcome?

Coin Flipping in Prolog

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```
p_coin :- heads.  
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What's the outcome?

```
?- coin.  
heads
```

```
?- p_coin.  
heads ; y
```

Coin Flipping in Prolog

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Prolog

```
p_coin :- heads.  
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```

What's the outcome?

```
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```
?- p_coin.  
heads ; y  
  
tails
```

- ▶ Prolog: backtracking
- ▶ CHR: **committed choice** for simplification rules

Propagation Coin Flipping

```
coin ==> heads.  
coin ==> tails.
```

What's the outcome?

```
?- coin.  
heads  
tails  
coin
```

Propagation Coin Flipping

```
coin ==> heads.  
coin ==> tails.
```

What's the outcome?

```
?- coin.  
heads  
tails  
coin
```

Rules are applied in sequence.

Propagation Coin Flipping

```
coin ==> heads.  
coin <=> tails.  
coin ==> side.
```

What's the outcome?

```
?- coin.  
heads  
tails
```

Propagation Coin Flipping

```
coin ==> heads.  
coin <=> tails.  
coin ==> side.
```

What's the outcome?

```
?- coin.  
heads  
tails
```

Rules are applied in sequence, until the active constraint is removed.

Exploiting Rule Order

```
generate(0) <=> true.  
generate(N) <=> N > 0 | value(N),  
                M is N - 1, generate(M).
```

Exploiting Rule Order

```
generate(0) <=> true.  
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vs.

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for queries ?- generate(N). with $N \geq 0$

Exploiting Rule Order

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generate(0) <=> true.  
generate(N) <=> N > 0 | value(N),  
                    M is N - 1, generate(M).
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vs.

```
generate(0) <=> true.  
generate(N) <=> value(N),  
                    M is N - 1, generate(M).
```

for queries ?- generate(N). with $N \geq 0$

Better Style?

Exploiting Rule Order

```
generate(0) <=> true.  
generate(N) <=> N > 0 | value(N),  
                M is N - 1, generate(M).
```

vs.

```
generate(0) <=> true.  
generate(N) <=> value(N),  
                M is N - 1, generate(M).
```

for queries ?- generate(N). with $N \geq 0$

Better Style?

Compare to:

```
generate(0) :- !.  
generate(N) :- value(N), M is N - 1, generate(M).
```

Rule Order

- ▶ rules are tried from top to bottom
- ▶ applied in sequence
- ▶ until the active constraint is removed
- ▶ committed choice: no alternatives explored for simplification rules



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- 10 Reactivation**
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Hello, World!

```
c(hello) <=> writeln(hello).  
c(world) <=> writeln(world).
```

```
?- c(hello), c(world).
```

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hello  
world
```

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```

```
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hello  
world  
?- c(X).
```

Hello, World!

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```

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?- c(hello), c(world).  
hello  
world  
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c(X)
```

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c(X)  
?- c(X), X = hello.  
hello  
?- c(X), c(world), X = hello.  
world  
hello  
X = hello
```

Reactivation

- ▶ constraints suspend in the constraint store
- ▶ unification reactivates suspended constraints

A Constraint: Inequality

```
neq(X,X) <=> fail.  
neq(X,Y) <=> X \= Y | true.
```

```
?- neq(a,a).
```

A Constraint: Inequality

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```
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neq(A,B)

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```
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```

neq(A,B)

```
?- neq(A,B), A = B.
```

No

```
?- neq(A,B), A = a, B = a.
```

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```

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```
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```

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```
?- neq(A,B), A = a, B = b.
```

Yes

```
?- neq(A,B), A = f(C), B = f(D).
```

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```

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```
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```
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```

Yes

```
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```

neq(f(C),f(D))

Using Reactivation for Finite Domains

```
domain(X,L)    <=> ground(X) | memberchk(X,L).
domain(X,[])   <=> fail.
domain(X,[V])  <=> X = V.
domain(X,L1), domain(X,L2) <=>
    intersection(L1,L2,L), domain(X,L).
intersection(L1,L2,L3) :- ...
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intersection(L1,L2,L3) :- ...
```

```
?- domain(X, [dog,fox,horse,snails,zebra]).
domain(X, [dog,fox,horse,snail,zebra])
?- domain(X, [d,f,h,s,z]), X = z.
```

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?- domain(X,[dog,fox,horse,snails,zebra]).
domain(X,[dog,fox,horse,snail,zebra])
?- domain(X,[d,f,h,s,z]), X = z.
X = z
?- domain(X,[d,f,h,s,z]), domain(X,[c,s,z]).
```

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X = z
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domain(X,[s,z])
```

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X = z
?- domain(X,[d,f,h,s,z]), domain(X,[c,s,z]).
domain(X,[s,z])
?- domain(X,[d,f,h,s,z]), domain(X,[c,z]).
```

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X = z
?- domain(X,[d,f,h,s,z]), domain(X,[c,s,z]).
domain(X,[s,z])
?- domain(X,[d,f,h,s,z]), domain(X,[c,z]).
X = z
```

Reactivation

- ▶ rules may not fire because of lack of instantiation

Reactivation

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- ▶ upon instantiation, the rule may now fire

Reactivation

- ▶ rules may not fire because of lack of instantiation
- ▶ upon instantiation, the rule may now fire
- ▶ useful for implementing constraint solvers



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Summary: CHR vs Prolog

	Prolog	CHR
heads	1	≥ 1
rule selection	unification	matching & guard
different rules	alternatives/backtracking	try all in sequence
no rule	failure	delay

CHR and Backtracking

```
p :- a.          a <=> c1.          b <=> d1.  
p :- b.          a <=> c2.          b <=> d2.
```

```
?- p.
```

```
c1
```

CHR and Backtracking

```
p :- a.          a <=> c1.          b <=> d1.  
p :- b.          a <=> c2.          b <=> d2.
```

```
?- p.  
c1 ; y  
d1
```

CHR and Backtracking

```
p :- a.          a <=> c1.          b <=> d1.  
p :- b.          a <=> c2.          b <=> d2.
```

```
?- p.  
c1 ; y  
d1
```

- ▶ Prolog creates choicepoints
- ▶ CHR does not
- ▶ Prolog backtracking undoes CHR changes

Backtracking for Labeling

```
domain(X,L)    <=> ground(X) | memberchk(X,L).
domain(X,[])   <=> fail.
domain(X,[V])  <=> X = V.
domain(X,L1), domain(X,L2) <=>
    intersection(L1,L2,L), domain(X,L).

split(X) <=> ground(X) | true.
split(X), domain(X,[V|Vs]) <=> X = V ; domain(X,Vs).
```

```
?- domain(X,[d,z]), split(X).
```


Backtracking for Labeling

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domain(X,L)    <=> ground(X) | memberchk(X,L).
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```

```
X = d ; y
```

```
X = z
```

Backtracking for Labeling

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domain(X,L1), domain(X,L2) <=>  
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```
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```

```
?- domain(X,[d,z]), split(X).
```

```
X = d ; y
```

```
X = z
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split(X) <=> ground(X) | true.
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```
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```
X = d ; y
```

```
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```

```
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```

```
X = d ; y
```

```
domain(X,[f,z])
```

Backtracking for Labeling

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```
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```

```
domain(X,L1), domain(X,L2) <=>  
    intersection(L1,L2,L), domain(X,L).
```

```
indomain(X) <=> ground(X) | true.
```

```
% indomain(X), domain(X,[V1,...,Vn]) <=> X = V1;...;X = Vn.
```

Backtracking for Labeling

```
domain(X,L)    <=> ground(X) | memberchk(X,L).
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% indomain(X), domain(X,[V1,...,Vn]) <=> X = V1;...;X = Vn.
indomain(X), domain(X,Vs) <=> member(X,Vs).
```

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?- domain(X,[d,f,z]), indomain(X).
```

```
X = d ; y
```

```
X = f ; y
```

```
X = z
```



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- ▶ how CHR works.
- ▶ how CHR differs from Prolog.

You should now have an understanding of:

- ▶ how CHR works.
- ▶ how CHR differs from Prolog.
- ▶ that CHR is not exclusively about constraints .



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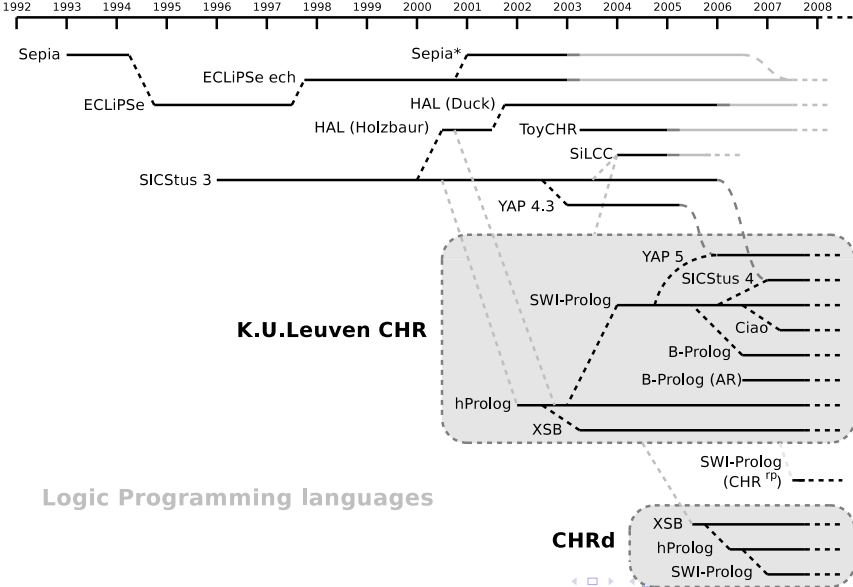
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- 2005- Leuven **JCHR** (Java), Peter Van Weert
- 2007 first **concurrent system**, Sulzmann & Lam

CHR Systems



Logic Programming languages

Thom Frühwirth, with his students:

- ▶ lots of example programs, constraint solvers and others

Thom Frühwirth & Tom Schrijvers,

- ▶ union-find

Jon Sneyers,

- ▶ Fibonacci heaps and Dijkstra's shortest path
- ▶ Hopcroft's DFA minimization
- ▶ Turing and RAM machine simulators

Henning Christiansen,

- ▶ meta-programming in CHR

...

Declarative Semantics

- ▶ classical first-order logic
- ▶ linear logic

(Frühwirth)

(Betz)

Declarative Semantics

- ▶ classical first-order logic (Frühwirth)
- ▶ linear logic (Betz)

Rewriting Properties

- ▶ confluence (Frühwirth; Duck; Haemmerle)
- ▶ completion (Frühwirth&Abdennadher)
- ▶ termination (Frühwirth; Voets&Pillozzi)
- ▶ complexity (Frühwirth; De Koninck)

Declarative Semantics

- ▶ classical first-order logic (Frühwirth)
- ▶ linear logic (Betz)

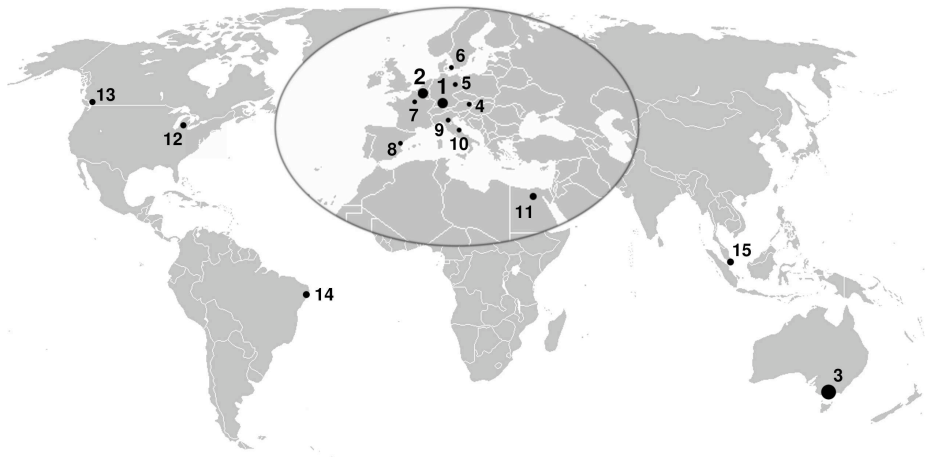
Rewriting Properties

- ▶ confluence (Frühwirth; Duck; Haemmerle)
- ▶ completion (Frühwirth & Abdennadher)
- ▶ termination (Frühwirth; Voets & Pillozzi)
- ▶ complexity (Frühwirth; De Koninck)

Applications

- ▶ type systems (Sulzmann & Stuckey)
- ▶ test case generation (Schrijvers)
- ▶ multi-agent systems (Alberti)
- ▶ ...

CHR Researchers around the World



- multi-headed business rules
- custom constraint solvers

- multi-headed business rules
- custom constraint solvers

▶ Scientific Software & Systems Ltd.

- ▶ stock brokering software



▶ Cornerstone Technology Inc

- ▶ injection mould design tool



▶ BSSE System and Software Engineering

- ▶ test generation



▶ MITRE Corporation

- ▶ optical network design



- ▶ CHR website (Google for)
 - ▶ programs
 - ▶ papers
- ▶ CHR Survey
 - ▶ As Time Goes By: Constraint Handling Rules – A Survey of CHR Research from 1998 to 2007, J. Sneyers, P. Van Weert, T. Schrijvers, L. De Koninck.

My Questions: Get the Quiz!

Your Questions?